



Radboud University



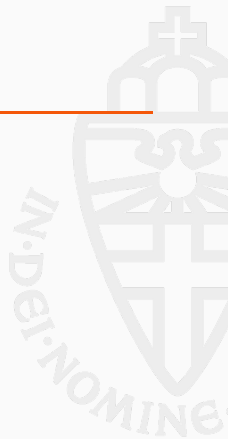
The transition to post-quantum crypto

Peter Schwabe

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<https://cryptojedi.org>

November 19, 2019



**DigiD**

Inloggen bij **Mijn DigiD**

Verplichte velden *

Inlogmethode *

- Ik wil inloggen met alleen gebruikersnaam en wachtwoord
- Ik wil inloggen met een controle via sms
- Ik wil inloggen met de DigiD app

DigiD gebruikersnaam ***Wachtwoord *** Onthoud mijn DigiD gebruikersnaam

U kunt tot 07:57 uur (Nederlandse tijd) inloggen. Daarna verloopt uw sessie.

Inloggen[Annuleren](#)[> Wachtwoord vergeten?](#)[> Nog geen DigiD? Vraag uw DigiD aan](#)



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Imiscweb1 DigiD: Inloggen | Inloggegevens - Vingerator

DigiD: Inloggen | Inlo... x +

https://digid.nl/inloggen

Veelgestelde vragen www.digid.nl

General Media Permissions Security

Website Identity

Website: **digid.nl**

Owner: **This website does not supply ownership information.**

Verified by: **KPN B.V.**

[View Certificate](#)

Privacy & History

Have I visited this website prior to today?	Yes, 9 times
Is this website storing information (cookies) on my computer?	Yes
Have I saved any passwords for this website?	No

[View Cookies](#)

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Technical Details

Connection Encrypted (TLS_ECDHE_RSA_WITH_AES_256_GCM_SHA384, 256 bit keys, TLS 1.2)

The page you are viewing was encrypted before being transmitted over the Internet. Encryption makes it difficult for unauthorized people to view information traveling between computers. It is therefore unlikely that anyone read this page as it traveled across the network.

[Help](#)

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Polynomial-Time Algorithms for Prime Factorization and Discrete Logarithms on a Quantum Computer*

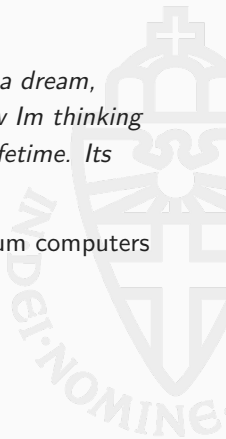
Peter W. Shor[†]

Abstract

A digital computer is generally believed to be an efficient universal computing device; that is, it is believed able to simulate any physical computing device with an increase in computation time by at most a polynomial factor. This may not be true when quantum mechanics is taken into consideration. This paper considers factoring integers and finding discrete logarithms, two problems which are generally thought to be hard on a classical computer and which have been used as the basis of several proposed cryptosystems. Efficient randomized algorithms are given for these two problems on a hypothetical quantum computer. These algorithms take a number of steps polynomial in the input size, e.g., the number of digits of the integer to be factored.

“In the past, people have said, maybe its 50 years away, its a dream, maybe itll happen sometime. I used to think it was 50. Now Im thinking like its 15 or a little more. Its within reach. Its within our lifetime. Its going to happen.”

—Mark Ketchen (IBM), Feb. 2012, about quantum computers



Definition

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5 main directions

- Lattice-based crypto (PKE and Sigs)
- Code-based crypto (mainly PKE)
- Multivariate-based crypto (mainly Sigs)
- Hash-based signatures (only Sigs)
- Isogeny-based crypto (so far, mainly PKE)



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- Inspired by two earlier NIST crypto competitions:
 - AES, running from 1997 to 2000
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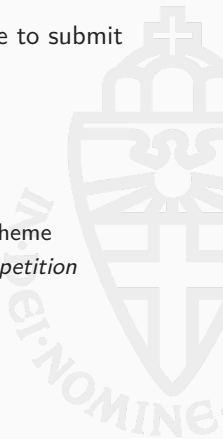
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- Selection through an open process and multiple rounds
- Actual decisions are being made by NIST



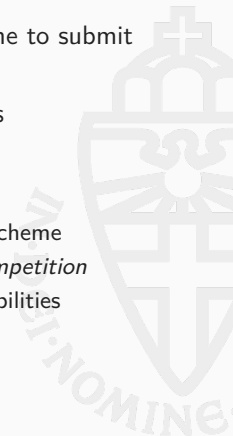
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 - Small tweaks are typically allowed, but standardized scheme represents state of the art *at the beginning of the competition*



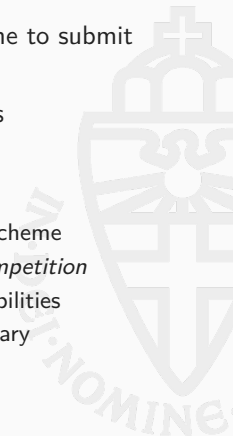
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- PQC project:
 - Announcement: Feb 2016
 - Call for proposals: Dec 2016 (based on community input)
 - Deadline for submissions: Nov 2017



The NIST competition, initial overview

Count of Problem Category	Column Labels		
Row Labels	Key Exchange	Signature	Grand Total
?	1		1
Braids	1	1	2
Chebychev	1		1
Codes	19	5	24
Finite Automata	1	1	2
Hash		4	4
Hypercomplex Numbers	1		1
Isogeny	1		1
Lattice	24	4	28
Mult. Var	6	7	13
Rand. walk	1		1
RSA	1	1	2
Grand Total	57	23	80

4 31 27

Overview tweeted by Jacob Alperin-Sheriff on Dec 4, 2017.

“Key exchange”

- What is meant is **key encapsulation mechanisms (KEMs)**
 - $(pk, sk) \leftarrow \text{KeyGen}()$
 - $(c, k) \leftarrow \text{Encaps}(pk)$
 - $k \leftarrow \text{Decaps}(c, sk)$

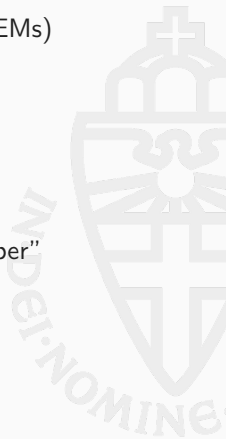


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Status of the NIST competition

- In total 69 submissions accepted as “complete and proper”
- Several broken, 5 withdrawn
- Jan 2019: NIST announces 26 round-2 candidates
 - 17 KEMs and PKEs
 - 9 signature schemes



Signature schemes

- 3 lattice-based
- 2 symmetric-crypto based
- 4 MQ-based



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KEMs/PKE

- 9 lattice-based
- 7 code-based
- 1 isogeny-based



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The baseline: ECC

- Today: build asymmetric crypto from elliptic-curve arithmetic
- Given P on a curve, $s \in \mathbb{Z}$, compute $Q = sP$
- ECDLP: hard to compute s , given P and Q
- Use for ECDH for key encapsulation and encryption
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- Performance (64-bit Intel CPU):
 - All operations between 50 000 and 200 000 cycles
 - Keys and ciphertexts: 32 bytes
 - Signatures: 64 bytes



PQ performance, some examples

- Supersingular-isogeny-based key agreement:
 - Public key/ciphertext: < 500 bytes each



PQ performance, some examples

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- McEliece code-based key agreement:
 - Encapsulation: $\approx 90\,000$ cycles
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 - Key generation: ≈ 300 Mio cycles
 - Cipher text: 188 bytes



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- \mathcal{MQ} -based signatures (e.g., GeMSS):
 - Signature: ≈ 50 bytes
 - Verification: $\approx 580\,000$ cycles
 - Signing: ≈ 2.7 billion cycles
 - Public key: ≈ 1.2 MB



Cryptographic hardness and proofs

- Need better understanding of attacks and their complexity
- Security reductions (“proofs”) help



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Secure implementations

- Implementations of secure schemes are not necessarily secure:
 - Subtle mistakes/bugs in implementations
 - Side-channel attacks
 - Fault attacks



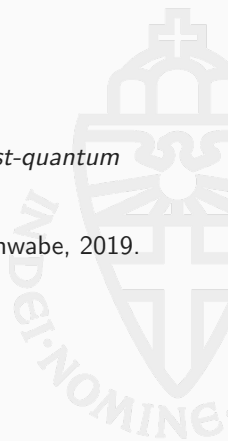
“the implementation security aspect of lattice-based cryptography is still a vastly unexplored and open topic”

—Primas, Pessl, Mangard, 2017.



“... this isn't very different for any of the other areas of post-quantum crypto”

—Schwabe, 2019.



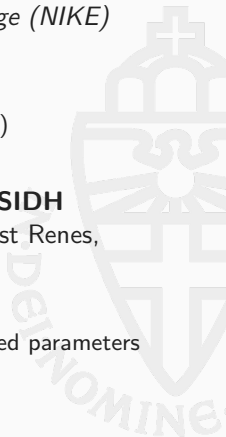
Challenges part 3: The case of DH

- Diffie-Hellman is extremely versatile:
- Can use it, for example, for *non-interactive key exchange (NIKE)*
 - Bob knows Alice' long-term public key A
 - Alice knows Bob's long-term public key B
 - They can each compute $k = h(A, B, aB) = h(A, B, bA)$
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- **Think protocols in KEMs, not in DHs/NIKEs!**

Challenges part 4: stateful signatures

- Hash-based signatures are already in RFCs:
 - XMSS: RFC8391
 - LMS: RFC8554
- Also highly parametrizable, for example:
 - Signing: ≈ 12.5 Mio cycles
 - Verification: ≈ 1 Mio cycles
 - Signature: ≈ 2.8 KB
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- XMSS gives forward security for free
- **Start thinking systems with *stateful* signatures**



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Shameless advertising

- pqm4: <https://github.com/mupq/pqm4>
- PQClean: <https://github.com/PQClean/PQClean>

