

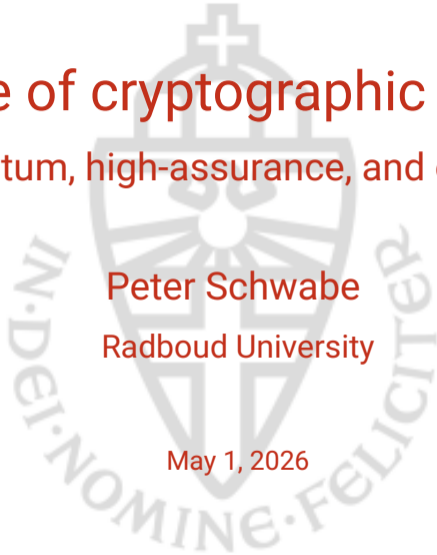
The Future of cryptographic engineering

post-quantum, high-assurance, and open-source

Peter Schwabe

Radboud University

May 1, 2026





*“An inaugural lecture is a formal public lecture given at a university by a **newly appointed** full professor, [...] It marks the professor’s official **introduction** to the academic community.”*

—Wikipedia on “Inaugural lecture” (emphasis added)



“Is this the full prof promotion from before Corona or another thing? Congrats in any case.”

—An answer to my invitation mail



1. Research motivated by **real-world challenges**
2. **Constructive** research: we like to build stuff
3. Scientific approach: Use **models and abstractions**, validate those in the lab

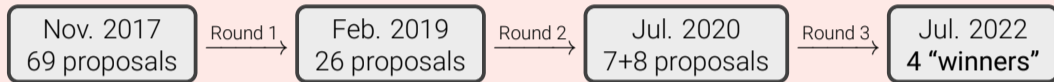
Polynomial-Time Algorithms for Prime Factorization and Discrete Logarithms on a Quantum Computer*

Peter W. Shor[†]

Abstract

A digital computer is generally believed to be an efficient universal computing device; that is, it is believed able to simulate any physical computing device with an increase in computation time by at most a polynomial factor. This may not be true when quantum mechanics is taken into consideration. This paper considers factoring integers and finding discrete logarithms, two problems which are generally thought to be hard on a classical computer and which have been used as the basis of several proposed cryptosystems. Efficient randomized algorithms are given for these two problems on a hypothetical quantum computer. These algorithms take a number of steps polynomial in the input size, e.g., the number of digits of the integer to be factored.

NIST PQC



"The public-key encryption and key-establishment algorithm that will be standardized is CRYSTALS-KYBER. The digital signatures that will be standardized are CRYSTALS-Dilithium, FALCON, and SPHINCS⁺. While there are multiple signature algorithms selected, NIST recommends CRYSTALS-Dilithium as the primary algorithm to be implemented"

—NIST IR 8413 (July 2022)

Where are we now?



Post-quantum encryption adoption

Post-quantum encrypted share of HTTPS request traffic ? 🔍 ⋮

Traffic type

Exclude bots



From <https://radar.cloudflare.com/adoption-and-usage#post-quantum-encryption-adoption>

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From <https://radar.cloudflare.com/adoption-and-usage#post-quantum-encryption-adoption>



So, what remains to be done in *research*?

crypto == pqc

EPOQUE: Engineering post-quantum cryptography

Secure implementations

- Optimize for tiny embedded systems and large Intel/AMD processors
- Optimize for speed and size, understand tradeoffs
- Side-channel attacks **and countermeasures**
- Fault attacks **and countermeasures**

Post-quantum protocols

- Re-think protocols with post-quantum primitives
- (How) can we live without Diffie-Hellman?
- Do we need efficient key generation?
- Do we need fast signatures?



Implementation security and **protocol migration**
are more relevant than ever!

IACR Transactions on Cryptographic Hardware and Embedded Systems

ISSN 2569-2925, Vol. 2023, No. 1, pp. 60–88.

DOI:10.46586/tches.v2023.i1.60-88

Adapting Belief Propagation to Counter Shuffling of NTTs

Julius Hermelink^{1,2}, Silvan Streit³,
Emanuele Strieder⁴ and Katharina Thieme⁵

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Ada **Belief Propagation Meets Lattice Reduction:
Security Estimates for Error-Tolerant Key
Recovery from Decryption Errors**

Julius Hermelink^{1†}, Erik Mårtensson^{2,3}, Simona Samardjiska⁴,
Peter Pessl⁵, Gabi Dreier Rodosek⁶

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Ada Se Bel The Insecurity of Masked Comparisons: SCAs on ML-KEM's FO-Transform

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Adaptability of Side-Channel Attacks on Masked Comparisons: SCAs on ML-KEM's EQ-Transform

A Generic Framework for Side-Channel Attacks against LWE-based Cryptosystems

Julius Hermelink¹, Silvan Streit^{2,3}, Erik Mårtensson^{4,5}, and Richard Petri¹

¹ Max Planck Institute for Security and Privacy, Bochum, Germany,
{julius.hermelink,richard.petri}@mpi-sp.org

² Fraunhofer AISEC, Garching, Germany, silvan.streit@aisec.fraunhofer.de

³ Technical University of Munich (TUM), Munich, Germany

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Ada Se Belin S.D. ...
**The Insecurity of Masked Comparisons: SCAs on ML-KEM's
EQ-Transform**

Max Pl **A Generic Framework for Side-Channel Attacks**

Ju

Max Pl

Julius I

**Finding and Protecting the Weakest Link
On Side-Channel Attacks on y in Masked ML-DSA**

² Fra

Julius Hermelink¹, Kai-Chun Ning¹, and Richard Petri¹

Max Planck Institute for Security and Privacy, Bochum, Germany,
firstname.lastname@mpi-sp.org

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² Fra

t-Probing (In-)Security

Pitfalls on Noise Assumptions

Dina Hesse¹, Jakob Feldtkeller^{2*}, Tim Güneysu¹, Julius Hermelink³,
Georg Land^{4*}, Markus Krausz⁵ and Jan Richter-Brockmann¹

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Julius I

² Fra

Noise-Tolerant Plaintext-Checking Oracle Attacks A Soft-Analytic Approach Applied to ML-KEM

Julius Hermelink¹, Erik Mårtensson^{2,3}, and Maggie Tran²

¹ Max Planck Institute for Security and Privacy, Bochum, Germany,
julius.hermelink@mpi-sp.org

² Lund University, Lund, Sweden, {erik.martensson,maggie.tran}@eit.lth.se

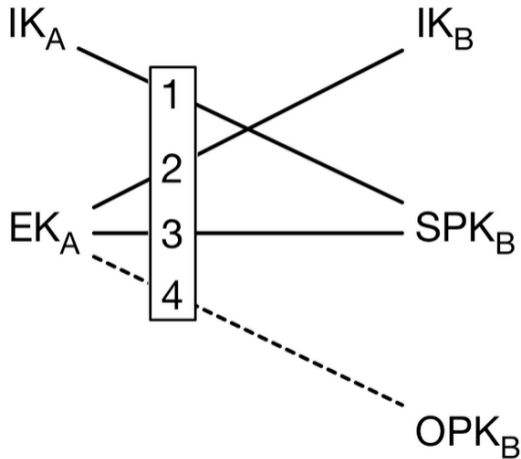
³ Advenica AB, Malmö, Sweden



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- ▶ Only goal: protect against HNDL

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- ▶ Original X3DH protocol

$$sk = \text{KDF}(\text{DH}_1|\text{DH}_2|\text{DH}_3|\text{DH}_4)$$

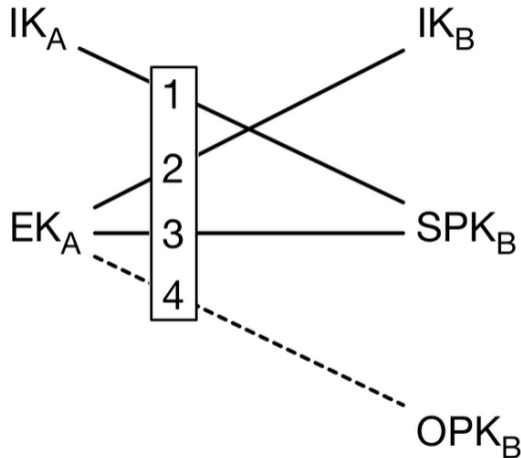


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 - ▶ Additionally obtain SS from PQ KEM
 - ▶ Compute final shared key as

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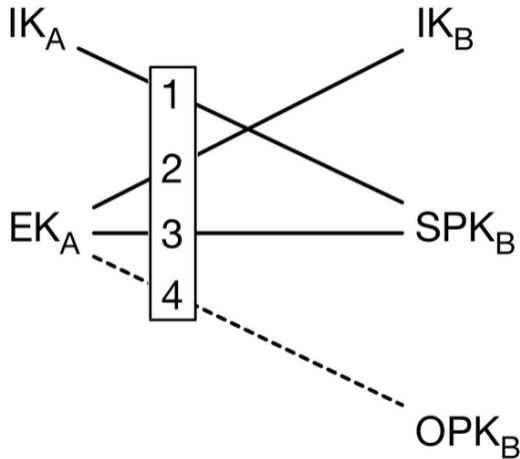
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$$sk = \text{KDF}(\text{DH}_1|\text{DH}_2|\text{DH}_3|\text{DH}_4|\text{SS})$$

- ▶ Two (potential) attacks discovered by Bhargavan, Jacomme, Kiefer, and Schmidt



More (KEMs) and signatures

- ▶ Research on block ciphers did not stop with FIPS-43-6 (DES)
- ▶ NIST PQC selected HQC as additional KEM in 2025
- ▶ NIST PQC continues selection of signatures (“on-ramp”)
- ▶ Can we get smaller, faster, more secure PQC?

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Advanced PQC

- ▶ KEMs and signatures are just a small part of public-key crypto
- ▶ PQ threshold schemes → NIST threshold competition
- ▶ PQ zero-knowledge and anonymous credentials



Challenges and Tools

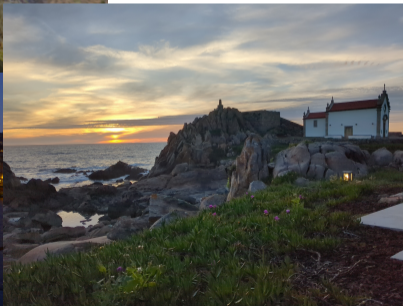
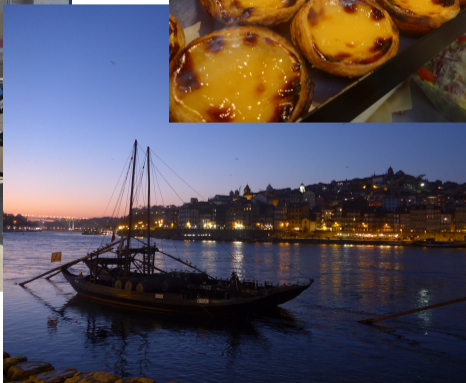
Challenges

- (Proven to be) secure implementations (including SCA security)
- Clear and meaningful security assumptions
- Verified security proofs using those assumptions
- Secure interfaces that connect to other components

Tools

- Today: Write software in C and assembly, ideally test/audit
- Need to move from 20th to 21st century:
 - Move from C to Rust (?)
 - Move from assembly to, e.g., jasmin (?)
 - Only use domain-specific languages (?)
 - Statically verify security properties during SW build

2020, the crazy year





Goal: Formally verified implementation of Kyber

≈9 papers, 30+ collaborators

Basavesh Ammanaghata Shivakumar,
Santiago Arranz Olmos, José Bacelar Almeida,
Gustavo Xavier Delerue Marinho Alves, Manuel
Barbosa, Francisca Barros, Gilles Barthe, Lionel
Blatter, Chitchanok Chuengsatiansup, Ignacio
Cuevas, François Dupressoir, Luís Esquível,
Ruben Gonzalez, Benjamin Grégoire, Andreas
Hülsing, Vincent Hwang, Jan Jancar, Matthias
Kannwischer, Vincent Laporte, Jean-Christophe
Léchenet, Ting-han Lim, Cameron Low, Tiago
Oliveira, Hugo Pacheco, Swarn Priya, Miguel
Quaresma, Rolfe Schmidt, Antoine Séré, Lucas
Tabary-Maujean, Pierre-Yves Strub, Yuval
Yarom, Zhiyuan Zhang, Jieyu Zheng

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FORMOSA CRYPTO

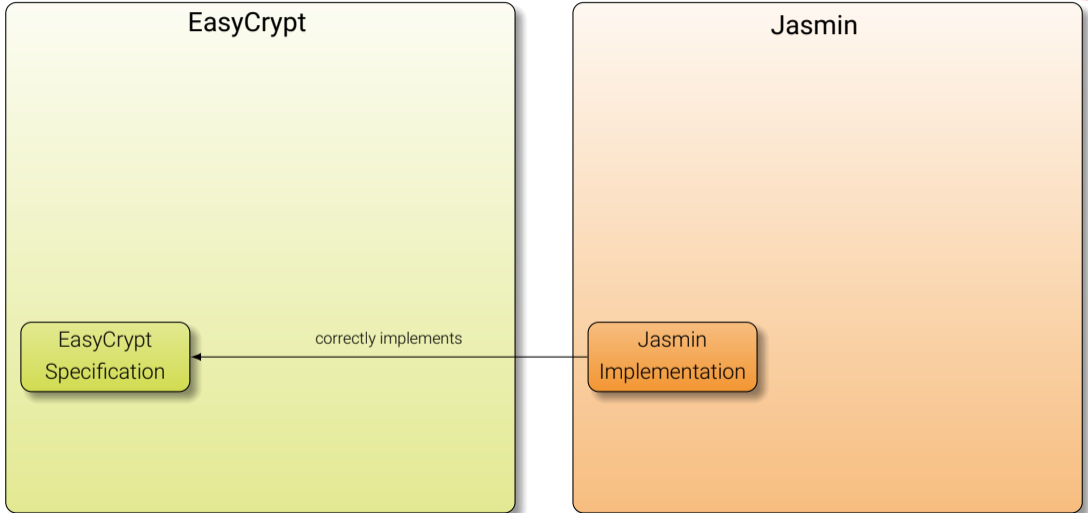


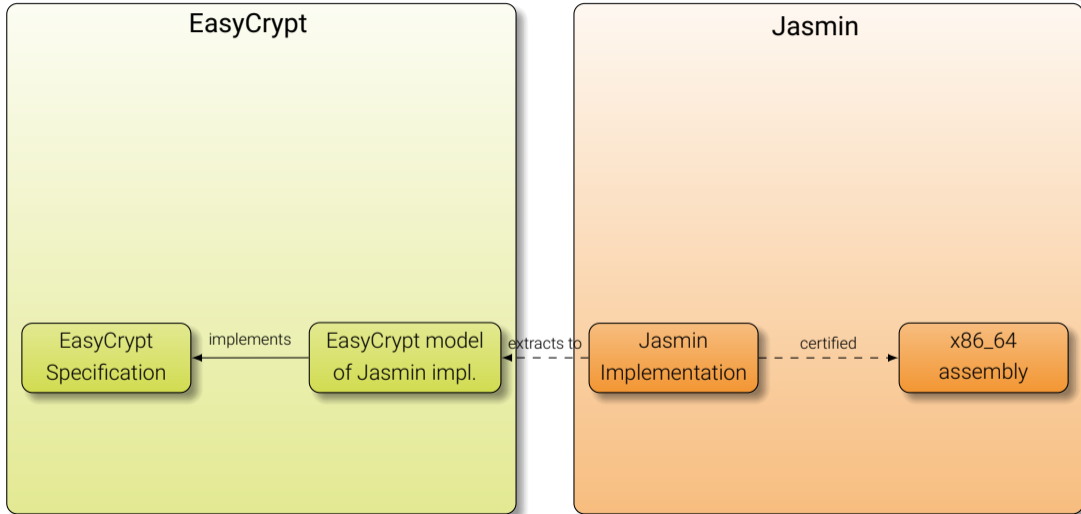
Digital Engineering - Universität Potsdam

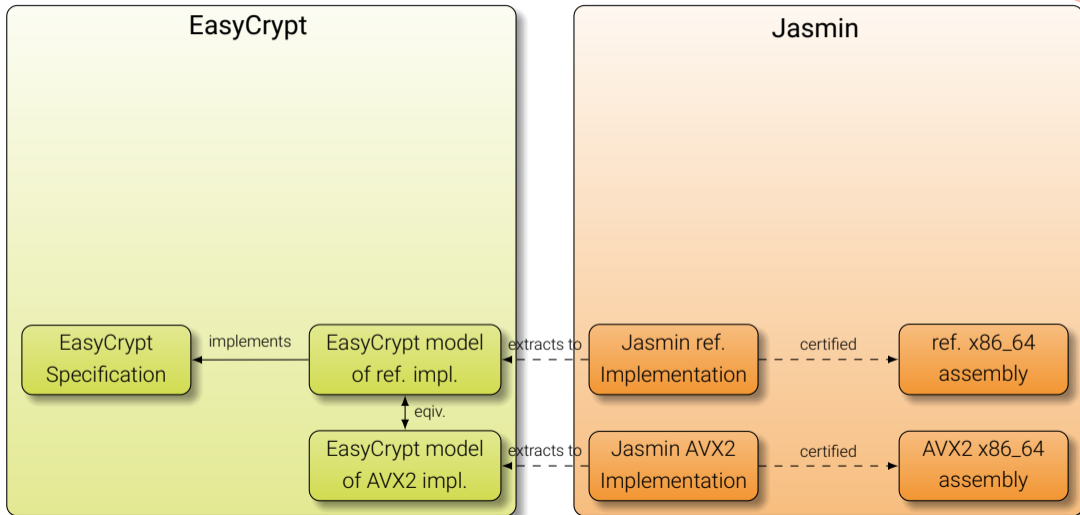


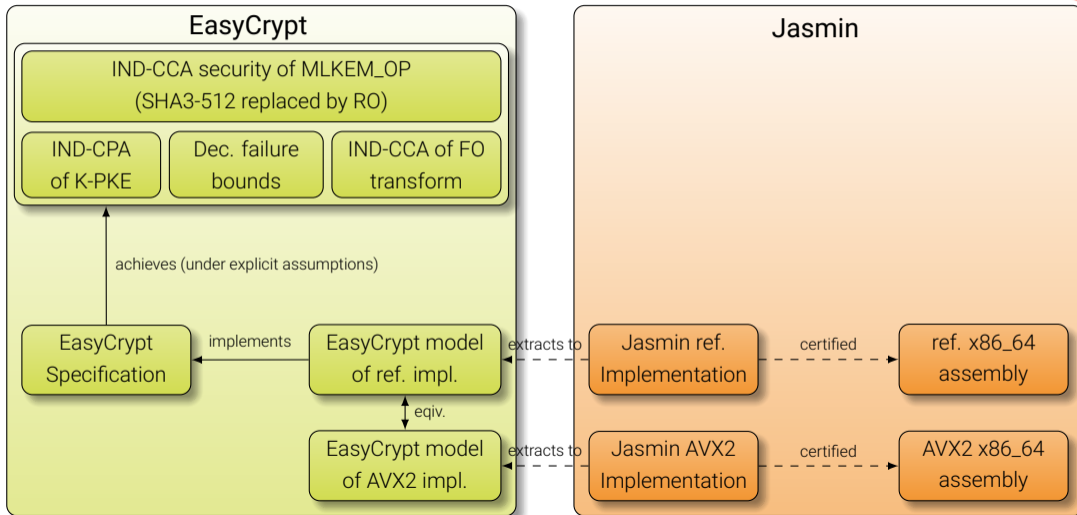
Universidade do Minho











Verifying Kyber

Hey Diego, we have a very important question. How hard would it be for us to change the title of our Formally verified Kyber paper at this point?

Specifically, change "Part 1" to "Episode IV"



Also, greetings from Lisbon, also from Manuel!

18:44

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Omg

That's the coolest and silliest request ever. Let me talk to Marcel but it should be fine

18:45

- ▶ **Future proof “constant-time”**
 - ▶ Information-flow type system in Jasmin
 - ▶ Only use DOIT instructions on secret data
 - ▶ Influenced update to DOIT instructions in Feb. 2026

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- ▶ **Systematic protection against Spectre attacks**

- ▶ Only DOIT on secret data also during transient execution
- ▶ Rewrite function returns to address Spectre-RSB
- ▶ Performance overhead of $<10\%$

▶ Future proof “constant-time”

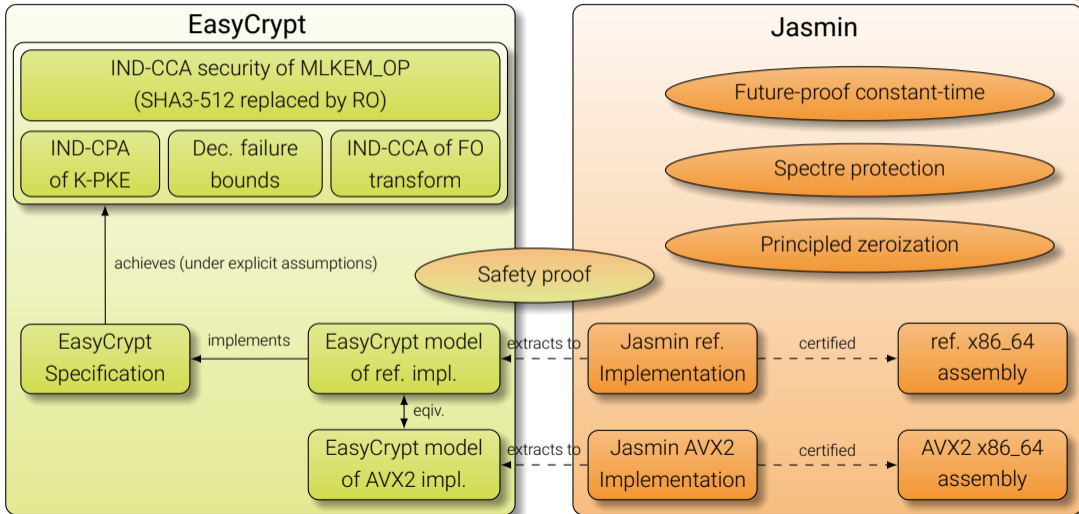
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- ▶ Performance overhead of <10%

▶ Principled zeroization

- ▶ Erase all (sensitive) data from stack and registers
- ▶ Automatic by Jasmin compiler
- ▶ Leverages various aspects of Jasmin



"We started using Kyber in production last year with our updated handshake protocol [...] Kyber is becoming an important building block for us.

Currently we rely on the PQClean implementation [...]

We are trying to decide what Kyber/ML-KEM implementation to use going forward. [...]

Is there anything that you think we should consider as we evaluate our options?

—Rolfe Schmidt, July 2024

Many fun challenges

- ▶ Consider more powerful attackers
- ▶ Super-optimization for Jasmin (CryptOpt, SLOTHY...)
- ▶ Link primitive and protocol proofs
- ▶ Go from protocols to applications

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One elephant in the room: **Scale!**

- ▶ Cannot take 6 years per primitive
- ▶ Already massively improved tooling
- ▶ Protocols more complex
- ▶ Applications *even* more complex

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One elephant in the room: **Scale!**

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- ▶ Protocols more complex
- ▶ Applications *even* more complex
- ▶ **GenAI to the rescue?**

What I didn't see coming in 2019: open source



Hang on, open source?

"It is very important to me that all my results can be freely used by anyone without any restrictions. I therefore place all my software in the public domain and make it available online."

—My Veni proposal (2013)

"Since the beginning of his academic career, the PI makes all results of his research, including software, freely available. All software on the PI's website <https://cryptojedi.org> is in the public domain (if requested via the Creative Commons Zero Waiver) to maximize reusability of results."

—My ERC proposal (2017)

“The fault attacker considered in this paper is thus limited to injecting only one fault per scalar multiplication. We assume that this fault may either skip a short block of consecutive instructions, set an arbitrary register to zero, or set an arbitrary register value to a random value [...] we also exclude the instruction pointer from the set of “arbitrary” registers that the attacker may fault.”

—Batina, Chmielewski, Haase, Samwel, Schwabe, 2023.

- ▶ Some attacks cannot be addressed in software (alone)
- ▶ Hardware we can work with does not have suitable protections
- ▶ Hardware with suitable protections is **very** closed

The frustration of a software guy



SCA-secure ECC in software – mission impossible?

SCA-secure ECC in software – mission impossible?

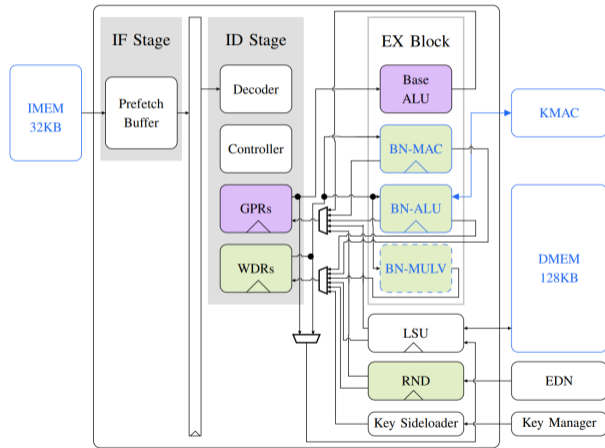
Mission impossible.

- ▶ RISC-V: open-source ISA
 - ▶ Started in 2010 at UC Berkeley
 - ▶ Multiple open-source implementations
 - ▶ Multiple commercial chips, now widely deployed
 - ▶ Security-minded design, but not purpose-built for crypto

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 - ▶ Security-minded design, but not purpose-built for crypto
- ▶ OpenTitan
 - ▶ Public launch in 2019
 - ▶ Open-source HW root of trust
 - ▶ Involvement from academia and industry
 - ▶ **Purpose-built for crypto**

Migrating open-source silicon to PQC

- ▶ Start with PQ secure boot (SLH-DSA)
- ▶ HW/SW support ML-KEM & ML-DSA
- ▶ Fork of OpenTitan
- ▶ OTBN → ACC
- ▶ Small area increase, massive speed boost
- ▶ Preparing industry and research tapeouts



Industry Partners

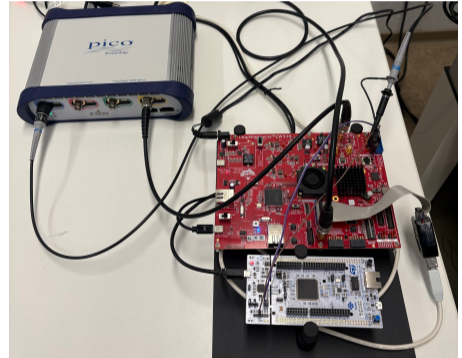


PQC as a catalyst for open-source silicon



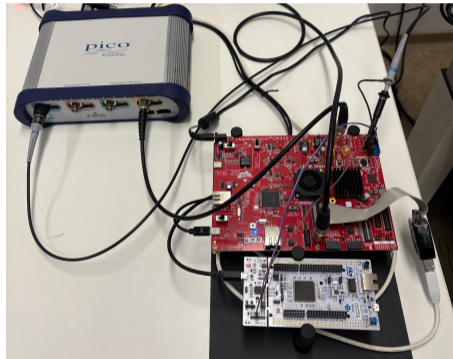
More powerful attacker: From software SCA to hardware SCA

- ▶ What can/should HW do for SW?
- ▶ What can/should SW do for HW?



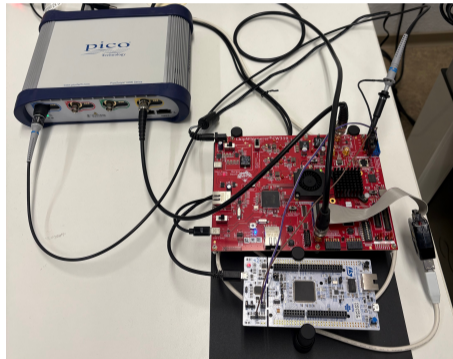
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 - ▶ What kind of noise actually matters?
 - ▶ Cost for increasing that noise (in HW and/or SW)
 - ▶ Cost for amplifying noise (masking order)
- ▶ Similar questions for fault injection. . .



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Goal: End-to-end tooling for principled SCA/FI security on open-source silicon

Think

- ▶ ML-KEM, ML-DSA, and post-quantum anonymous credentials,

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- ▶ implemented in Jasmin,

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- ▶ offering well-defined interfaces to, e.g., Signal and PQ-eIDAS, and

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- ▶ proofs of security extending to these protocols and applications.

A massive team effort



epoqc

<https://epoqc.org>

stay tuned...