## Hacking in C Pointers

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# Allocation of multiple variables

Consider the program

main(){
 char x;
 int i;
 short s;
 char y;
 ....
}

What will the layout of this data in memory be? Assuming 4-byte ints, 2-byte shorts, and little endian architecture

#### Printing addresses where data is located

We can use & to see where data is located

```
char x; int i; short s; char y;
```

```
printf("x is allocated at %p \n", &x);
printf("i is allocated at %p \n", &i);
printf("s is allocated at %p \n", &s);
printf("y is allocated at %p \n", &y);
    // Here %p is used to print pointer values
```

Compiling with or without -02 will reveal different alignment strategies

Memory as a sequence of bytes

	<b>x i</b> <sub>0</sub>	$\mathbf{i}_1$ $\mathbf{i}_2$	<b>i</b> 3	<b>s</b> 0	<b>s</b> <sub>1</sub>	У			
--	-------------------------	-------------------------------	------------	------------	-----------------------	---	--	--	--

But on a 32-bit machine, the memory is a sequence of 4-byte words

х	<b>i</b> 0	$\mathbf{i}_1$	$\mathbf{i}_2$
<b>i</b> 3	<b>S</b> 0	<b>s</b> <sub>1</sub>	у

Now the data elements are not nicely aligned with the words, which will make execution slow, since CPU instructions act on words.

Different allocations, with better/worse alignment

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Lousy alignment, but uses minimal memory

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nt, Possible compromise

Optimal alignment, but wastes memory Compilers may introduce padding or change the order of data in memory to improve alignment.

There are trade-offs here between speed and memory usage.

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Most C compilers can provide many optional optimizations. E.g., use  $$\tt man\ gcc$$ 

to check out the many optimization options of gcc.

# Arrays

#### Arrays

An array contains a collection of data elements with the same type. The size is constant.

```
int test_array[10];
int a[] = {30,20};
test_array[0] = a[1];
```

```
printf("oops %d \n", a[2]); //will compile & run
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## Arrays

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#### Array bounds are **not** checked.

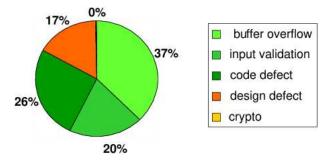
Anything may happen when accessing outside array bounds. The program may crash, usually with a segmentation fault (segfault).

The historic decision **not** to check array bounds is responsible for in the order of 50% of all the security vulnerabilities in software, in the form of so-called buffer overflow attacks.

Other languages took a different (more sensible?) choice here. E.g. ALGOL60, defined in 1960, already included array bound checks.

# Typical software security vulnerabilities

Security bugs found in Microsoft's first security bug fix month (2002)



Here buffer overflows are platform-specific.

Some of the *code defects* and *input validation* problems might also be. *Crypto* problems are much more rare, but can be of very high impact.

## Array bounds checking

Tony Hoare in Turing Award speech on the design principles of ALGOL 60

"The first principle was security: ... A consequence of this principle is that every subscript was checked at run time against both the upper and the lower declared bounds of the array. Many years later we asked our customers whether they wished us to provide an option to switch off these checks in the interests of efficiency. Unanimously, they urged us not to –



they knew how frequently subscript errors occur on production runs where failure to detect them could be disastrous. I note with fear and horror that even in 1980, language designers and users have not learned this lesson. In any respectable branch of engineering, failure to observe such elementary precautions would have long been against the law."

[ C.A.R.Hoare, The Emperor's Old Clothes, Communications of the ACM, 1980]

# Overrunning arrays

Consider the program

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char a[2];
int x = 6;
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By overrunning an array we can try to reverse-engineer the memory layout

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E.g. a compiler might allocate 16 bytes rather than 15 bytes for

```
char text[15];
```

#### Arrays are passed by reference

Arrays are always passed by reference.

```
For example, given the function
```

```
void increase_elt(int x[]) {
    x[1] = x[1]+23;
}
```

What is the value of a [1] after executing the following code?

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int a[2] = {1, 2};
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25

Recall call by reference from Imperative Programming

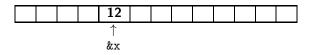
# Pointers

Retrieving addresses of *pointers* using &

We can find out where some data is allocated using the & operation. If

int x = 12;

then &x is the memory address where the value of x is stored, aka a pointer to x.

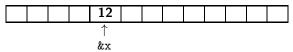


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It depends on the underlying architecture how many bytes are needed to represent addresses: 4 on 32-bit machines, 8 on a 64-bit machine.

Pointers are typed:

the compiler keeps track of what data type a pointer points to

int \*p; // p is a pointer that points to an int
float \*f; // f is a pointer that points to a float

Suppose

```
int y, z; int *p; // i.e. p points to an int
```

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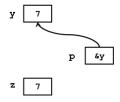
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How can we get the value that a pointer points to? Using \*

Looking up what a pointer points to, with \*, is called dereferencing.

# Confused? draw pictures!



What is the value of y?

int y = 2; int x = y; y++; x++;

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What is the value of y?

int y = 2; int \*x = &y; y++; (\*x)++; Note that \* is used for 3 different purposes, with 3 different meanings

2. As a prefix operator on pointers

int z = \*p;

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Some legal C code can get confusing, e.g.

z = 3 \* \*p

What can be confusing in

int \*p = &y;

is that this is an assignment to p, not \*p

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Downside of writing int\*

int\* x, y, z;

declares x as a pointer to an int and y and z as int...

# Still not confused?

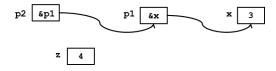
What will the value of z be?

## Still not confused?

What will the value of z be?

What should the types of p1 and p2 be?

## Still not confused? pointers to pointers



What is the value of y at the end?

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```
What is the value of y at the end? 6
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```

```
What is the value of y at the end?
6
What is the value of *p at the end?
6
What is the value of *g at the end?
```

```
int y = 2;
int z = 3;
int* p = &y;
int* q = &z;
(*q)++;
*p = *p + *q;
q = q + 1;
printf("y is %d\n", y);
```

```
What is the value of y at the end?
6
What is the value of *p at the end?
6
What is the value of *q at the end?
We don't know! q points to some memory cell after z in the memory
```

## Pointer arithmetic

You can use + and - with pointers.

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```

```
For example, if
     int *ptr; char *str;
then
    ptr + 2 means "Add 2 * sizeof(int) to the address in ptr"
    str + 2 means "Add 2 to the address in str"
(because sizeof(char) is 1)
```

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Suppose

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Suppose

Now

p + 3

points to

a[3]

So we use addition to pointer p to move through the array

## Pointer arithmetic for strings

What is the output of

```
char *msg = "hello world";
char *t = msg + 6;
printf("t points to the string %s.", t);
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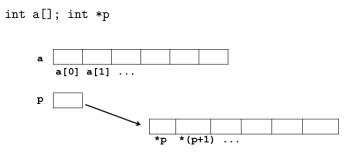
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This will print

t points to the string world.

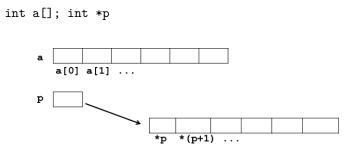
## Arrays vs pointers

Arrays and pointers behave similarly, but are very different in memory Consider



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A difference: a will always refer to the same array, whereas p can point to different arrays over time

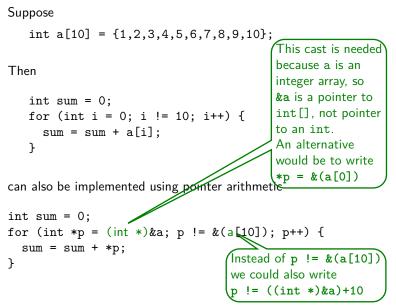
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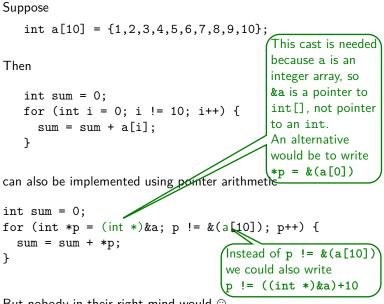
int a[10] = {1,2,3,4,5,6,7,8,9,10};

Then

```
int sum = 0;
for (int i = 0; i != 10; i++) {
  sum = sum + a[i];
}
```

can also be implemented using pointer arithmetic





But nobody in their right mind would ©

A problem with pointers: ...

```
int i; int j; int *x;
...
// lots of code omitted
i = 5;
j++
// what is the value of i here?
(*x)++;
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i = 5;
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(*x)++;
// what is the value of i here? 5 or 6, depending on
whether *x points to
i
```

Two pointers are called **aliases** if they point to the same location

```
int i = 5;
int *x = &i;
int *y = &i;
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(*x)++;
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Keeping track of pointers, in the presence of potential aliasing, can be really confusing, and really hard to debug...

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Confusingly, the meaning of addition for pointers depends on their type, as +1 for pointers of type int \* really means +sizeof(int)

The potential of pointers: inspecting raw memory

To inspect a piece of raw memory, we can cast it to a

```
unsigned char *
```

and then inspect the bytes

```
float f = 3.14;
unsigned char *p = (unsigned char *) &f;
printf("The representation of float %f is", f);
for (int i = 0; i < sizeof(float); i++, p++);) {
printf("%d", *p);
}
printf("\n");
```

intptr\_t defined in stdint.h is an integral type that is guaranteed to be wide enough to hold pointers.

```
int *p; // p points to an int
intptr_t i = (intptr_t) p; // the address as a number
p++;
i++;
// Will i and p be the same?
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There is also an unsigned version of intptr\_t: uintptr\_t

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As other arrays, we can use both the array type char [] and the pointer type char \* for them.



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Working with C strings is highly error prone! There are two problems

- 1. As for any array, there are no array bounds checks so it's the programmer's responsibility not to go outside the array bounds
- It is also the programmer's responsibility to make sure that the string is properly terminated with a null character.
   If a string lacks its null terminator, e.g. due to problem 1, then standard functions to manipulate strings will go off the rails.

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Moral of the story: if you can, avoid using standard C strings. E.g. in C++, use C++ type strings; in C, use safer string libraries.

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Compilers can emit warnings if you change string literals, e.g.

```
gcc -Wwrite-strings
```

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- Their representation
- How these representations can be 'broken', i.e. how we can inspect and manipulate the underlying representation (e.g. with casts)
- Some things that can go wrong e.g. due to access outside array bounds or integer under/overflow